

A Framework for Fairness in Two-Sided Marketplaces

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Outline

- Introduction
- Background
- Two-Sided Fair Marketplace
- Experiments
- Conclusion

Motivation



Job Seeker Example

Singh, Ashudeep, and Thorsten Joachims. "Fairness of exposure in rankings." *Proceedings of the 24th ACM SIGKDD International Conference on Knowledge Discovery & Data Mining*. 2018.

Suggestions: one gender!

Search for CEO in a serach engine

Introduction

Recommender systems

Single-sided

members viewing:

- Items
- Products
- Jobs
- Movies
- Restaurants

Two-sided member-to-member marketplace

when members can serve multiple functions:

- feed ranking
- people or friend recommendations
- search systems
- recruiters searching for job candidates

Introduction Goal:

Guarantee that the rankings are fair to:

- The source members initiating the queries
- The destination members who are being returned by the query

Homś

- Define a multi-session utility and add fairness constraints for destination members across multiple sessions
- Add fairness constraints for source members
- Scale for large-scale recommender systems by using a duality idea

People You May Know (Linkedin)

- PYMK is a Member-Member Marketplace
- Fairness based on Gender attribute (Similar setting for IM, FM)
- Destination Side Fairness Directly observable by members
- Source Side Fairness Unobservable by individual members
 - Can be exposed if audited

Source: Member who looks at ranking

Suggestions: one gender!



Destination: Members being ranked

A General Framework

Fairness of Exposure in Rankings

Ashudeep Singh Cornell University Ithaca, NY ashudeep@cs.cornell.edu							
ABSTI	$\mathbf{P} = \operatorname{argmax}_{\mathbf{P}} \mathbf{u}^T \mathbf{P} \mathbf{v}$						
Ranking							
transitio	$\mathbf{r} + 1^T \mathbf{p} - 1^T $						
jobs, job	S.I. $\square \square \square$ (Sum						
is a subst							
not only	$\mathbf{P} = \mathbf{I} (sum c)$						
not only							
is a subst	$0 < \mathbf{P}_{i,i} < 1$						
jobs, job	$V _ I l, J _ I$						
transitio	P is fair						

Thorsten Joachims Cornell University Ithaca, NY tj@cs.cornell.edu

(expected utility) ob seek-

n of probabilities for each position)

of probabilities for each document)

(valid probability)

(fairness constraints)

ob seekne of the king sysobability g should elevance,

sing sysbability g should levance,



Preliminaries

Probabilistic Rankings

P_{i,j}: the probability that ranking R places item **i** at rank **j**

Birkhoff-von Neumann (BvN) decomposition:

doubly stochastic matrix into a convex sum of permutation matrices

$\mathbf{A} = \theta_1 \mathbf{A}_1 + \theta_2 \mathbf{A}_2 + \dots + \theta_n \mathbf{A}_n$

P forms a doubly stochastic matrix of size N × N

$$P = \begin{bmatrix} 0.4 & 0.5 & 0.1 \\ 0.3 & 0.3 & 0.4 \\ 0.3 & 0.2 & 0.5 \end{bmatrix}$$

$$\mathbf{R}' = \begin{vmatrix} 0 & 0.4 & 0.2 & 0.4 \\ 0.4 & 0.4 & 0 & 0.2 \\ 0.4 & 0.2 & 0.4 & 0 \\ 0.2 & 0 & 0.4 & 0.4 \end{vmatrix} = \mathbf{0.4} \begin{vmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{vmatrix} + \mathbf{0.4} \begin{vmatrix} 0 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{vmatrix}$$
$$+ \mathbf{0.2} \begin{vmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{vmatrix}$$



Probabilistic Rankings

Expected utility for a probabilistic ranking:

Rewrite it as a matrix product:

Discounted Cumulative Gain (DCG) can be represented in this format: P is a doubly stochastic matrix of size N × N



$$U(\mathbf{P}|q) = \sum_{d_i \in \mathcal{D}} \sum_{j=1}^{N} \mathbf{P}_{i,j} u(d_i|q) v(j)$$

$$\mathbf{U}(\mathbf{P}|q) = \mathbf{u}^T \mathbf{P} \mathbf{v}$$

$$DCG(r|q) = \sum_{u \in \mathcal{U}} P(u|q) \sum_{d \in \mathcal{D}} \frac{2^{\operatorname{rel}(d|u,q)} - 1}{\log(1 + \operatorname{rank}(d|r))}$$



Optimizing Fair Rankings via Linear Programming

- maximizing source side utility
- Solve problem per session!

P(i, j): probability of showing the i-th destination member in the j-th slot

 $\mathbf{P} = \operatorname{argmax}_{\mathbf{P}} \ \mathbf{u}^{T} \mathbf{P} \mathbf{v}$ s.t. $\mathbb{1}^{T} \mathbf{P} = \mathbb{1}^{T}$ (sum of $\mathbf{P}\mathbb{1} = \mathbb{1}$ (sum of $0 \le \mathbf{P}_{i,j} \le 1$ \mathbf{P} is fair

Singh, Ashudeep, and Thorsten Joachims. "Fairness of exposure in rankings." *Proceedings of the 24th ACM SIGKDD International Conference on Knowledge Discovery & Data Mining*. 2018.

fairness constraints on the destination members.

(expected utility)

s.t. $\mathbb{1}^T \mathbf{P} = \mathbb{1}^T$ (sum of probabilities for each position)

(sum of probabilities for each document)

(valid probability)

(fairness constraints)

Fairness Measures

 Destination Side Metrics • **Demographic Parity** (Equality of Exposure) Average exposure of the groups are equal in each query.

• Disparate Treatment

Average exposure of the groups proportional to their utility are equal in each query.

• Disparate Impact query

$Exposure(G_0|\mathbf{P}) = Exposure(G_1|\mathbf{P})$

Exposure($G_0 \mathbf{P}$)	Exposure($G_1 \mathbf{P}$)		
$U(G_0 q)$	$ U(G_1 q)$		

Average impact (eg CTR) of the groups proportional to their utility are equal in





Position Bias
Demographic Parity
Avg Exposure for a group
Exposure(
$$G_k | \mathbf{P} \rangle = \frac{1}{|G_k|} \sum_{d_i \in G_k} \text{Exposure}(d_i | \mathbf{P})$$

Average exposure of the groups are equal in each query.
Exposure($G_0 | \mathbf{P} \rangle = \text{Exposure}(G_1 | \mathbf{P})$
 $\Leftrightarrow \frac{1}{|G_0|} \sum_{d_i \in G_0} \sum_{j=1}^N \mathbf{P}_{i,j} \mathbf{v}_j = \frac{1}{|G_1|} \sum_{d_i \in G_1} \sum_{j=1}^N \mathbf{P}_{i,j} \mathbf{v}_j$
 $\Leftrightarrow \sum_{d_i \in \mathcal{D}} \sum_{j=1}^N \left(\frac{1}{d_i \in G_0} - \frac{1}{|G_1|} \right) \mathbf{P}_{i,j} \mathbf{v}_j = 0$
 $\Leftrightarrow \mathbf{f}^T P \mathbf{v} = 0$ (with $\mathbf{f}_i = \frac{1}{|G_0|} - \frac{1}{|G_1|} \right)$

Disparate Treatment

Average exposure of the grou equal in each query.





Average exposure of the groups proportional to their utility are

$$\frac{\mathbf{P}}{\mathbf{P}} = \frac{\text{Exposure}(G_1|\mathbf{P})}{\mathbf{U}(G_1|q)} \\
\frac{\mathbf{D}_{j=1}^{N} \mathbf{P}_{i,j} \mathbf{v}_j}{\mathbf{Q}} = \frac{\frac{1}{|G_1|} \sum_{d_i \in G_1} \sum_{j=1}^{N} \mathbf{P}_{i,j} \mathbf{v}_j}{\mathbf{U}(G_1|q)} \\
\frac{\mathbf{P}_{i,j} \in G_0}{\mathbf{U}(G_0|q)} - \frac{\frac{1}{d_i \in G_1}}{|G_1|\mathbf{U}(G_1|q)} \mathbf{P}_{i,j} \mathbf{v}_j = 0 \\$$
(with $\mathbf{f}_i = \left(\frac{1_{d_i \in G_0}}{|G_0|\mathbf{U}(G_0|q)} - \frac{1_{d_i \in G_1}}{|G_1|\mathbf{U}(G_1|q)}\right)$





 $P(\text{click on document } i) = P(\text{examining } i) \times P(i \text{ is relevant})$ = Exposure($d_i | \mathbf{P}$) × P(i is relevant) $= \left(\sum_{i=1}^{I} \mathbf{P}_{i,j} \mathbf{v}_j\right) \times \mathbf{u}_i$

- $\frac{\operatorname{CTR}(G_0|\mathbf{P})}{\operatorname{U}(G_0|q)} = \frac{\operatorname{CTR}(G_1|\mathbf{P})}{\operatorname{U}(G_1|q)}$ $\frac{\frac{1}{|G_0|} \sum_{i \in G_0} \sum_{j=1}^{N} \mathbf{P}_{i,j} \mathbf{u}_i \mathbf{v}_j}{U(G_0|q)} = \frac{\frac{1}{|G_1|} \sum_{i \in G_1} \sum_{j=1}^{N} \mathbf{P}_{i,j} \mathbf{u}_i \mathbf{v}_j}{U(G_1|q)}$ $\left(\frac{\mathbb{1}_{d_i \in G_0}}{|G_0| U(G_0|q)} - \frac{\mathbb{1}_{d_i \in G_1}}{|G_1| U(G_1|q)}\right) \mathbf{u}_i \mathbf{P}_{i,j} \mathbf{v}_j = 0$ (with $\mathbf{f}_i = \left(\frac{\mathbbm{1}_{d_i \in G_0}}{|G_0| U(G_0|q)} - \frac{\mathbbm{1}_{d_i \in G_1}}{|G_1| U(G_1|q)} \right) \mathbf{u}_i \right)$



Two-sided Fair Marketplace

optimization-based framework

- maximizing source side utility
- maintaining fairness constraints on the destination members.
- post-processing
- Detailed Design

Maximize $u_s^{+} P_s v$ Subject to $1^{\top} P_s = 1$, $P_s 1 \le 1$, $0 \le P_s(d, r) \le 1$, $f_{k,k'}^{\top} P_s w = 0$ and

Expanded framework

- Fairness constraints for destination members at multiple sessions by considering the multi-session utility
- Adding fairness constraints for source members
- Post-processing
- Detailed Design

Destination Side Fairness

Multi-Session Fairness

 $\tilde{u}_s P_s w = c$

Two-Sided Fair Marketplace

A Post-Processesing approach:

- Operates on the final ranking with fairness constraints (very generally applicable)
- Agnostic to the underlying model architecture
- Suitable for large-scale recommender systems

Two-sided Fair Marketplace

Most recommender systems:

maximizing source side utility



P(d, r): probability of showing the d-th destination member in the r-th slot for source s and query q

Expanded Framework:

Destination side utility

- 1. much harder to tackle
- 2. depends of different queries arising from different sources



$${}^{e}(s) = \sum_{d=1}^{D_{q}} \sum_{r=1}^{m} u_{s,d}^{q} P_{s}^{q}(d,r) v_{r} = u_{s}^{\top}$$

$$(s,d) = u_{s,d}^q \sum_r P_s^q(d,r) \cdot v_r = u_{s,d}^q P_s^q(d,r)$$





Destination Side Utility

expected utility a destination member d receives

average cumulative utility for group G at time t



ρ is a



incremental utility



equality of average incremental utility across groups (Multi-session Dest-side Fairness)

$$e^{est}(d)[T] = \sum_{i=1}^{N} \rho^{T-t_i} U_{q_i}^{dest}(s_{t_i}, d)$$

In appropriately chosen discount rate
 $G_k[t] = \frac{1}{|G_k|} \sum_{d \in G_k} U^{dest}(d)[t]$

$$\mu_{G_k}[T] = \mu_{G_k}[T] - \mu_{G_k}[t]$$

$$\Delta \mu_{G_k}[T] = \Delta \mu_{G_{k'}}[T]$$



Multi-session Destination Side Fairness Constraint





$$] = \Delta \mu_{G_{k'}}[T]$$

Two-sided Fair Marketplace Optimization Details **Optimization problem with multi-session Fairness:** Maximize $u_s^{\top} P_s v$ Subject to $1^{\top}P_s = 1$, $P_s 1 \le 1$, $0 \le P_s$

Vectorized form:

$$\begin{aligned} & \operatorname{Max} \, p^\top (u \cdot v) \\ & \text{s.t.} \ \ p^\top (f_{k,k'} \cdot v) = 0, \ \ p^\top (\tilde{u} \cdot v) = c, \end{aligned}$$

Add a regularizer + better format

Maximize
$$p^{ op}(u \cdot v) + \frac{\gamma}{2}p^{ op}p$$
 sub

$$P_s(d,r) \leq 1, \quad egin{array}{c} ext{Destination Side Fairness} \ f_{k,k'}^ op P_s w = 0 \ f_{k,k'}^ op P_s w = 0 \ ext{and} \quad egin{array}{c} ext{Multi-Session Fairness} \ ilde{u}_s P_s w = c \ ilde{$$

$[I_m:I_m:\cdots:I_m]p=1$ and $p_d\in T_m$ $\forall d\in M$

bject to $Ap \leq b$ and $p_d \in T_m \quad \forall \ d \in M$



Two-sided Fair Marketplace Optimization Details

Maximize
$$p^{ op}(u \cdot v) + \frac{\gamma}{2}p^{ op}p$$
 sul

- Solve the optimization problem
- Obtain the dual variables $(\lambda_1, \lambda_2, \eta)$
- Obtain the primal solution from duals (Drastically reduces the latency cost)

$$\hat{p}_d(\lambda_1,\lambda_2,\eta) = \Pi_{T_m} \Big(rac{1}{\gamma} \left\{ (u \cdot v) \right\}$$

bject to $Ap \leq b$ and $p_d \in T_m \quad \forall \ d \in M$

 $_d - \lambda_1 (f_{k,k'} \cdot v)_d - \lambda_2 (\tilde{u} \cdot v)_d - \eta \}$

Extension to Source Side Fairness

multiple sessions.

$$\mathbb{E}\left[U^{\text{source}}(G_{k},T)\right] = \sum_{i \notin j \leq T} \rho^{T-t_{i}} U^{\text{source}}_{q_{i}}(s_{j}) = \sum_{i:t_{i} \leq T} \rho^{T-t_{i}} u^{\top}_{s_{i}} P_{s_{i}}$$
$$\mathbb{E}\left[U^{\text{source}}(G_{k},T+\delta)\right] = \mathbb{E}\left[U^{\text{source}}(G_{k'},T+\delta)\right]$$
$$u^{\top}_{s} P_{s} v = \mathbb{E}\left[U^{\text{source}}(G_{k'},T+\delta)\right] - \rho^{\delta} \mathbb{E}\left[U^{\text{source}}(G_{k'},T)\right]$$

$$\mathbb{E}\left[U^{\text{source}}(G_{k},T)\right] = \sum_{i \notin j \leq T} \rho^{T-t_{i}} U^{\text{source}}_{q_{i}}(s_{j}) = \sum_{i:t_{i} \leq T} \rho^{T-t_{i}} u^{\top}_{s_{i}} P_{s_{i}}$$
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$$\mathbb{E}\left[U^{\text{source}}(G_{k},T)\right] = \sum_{i \notin j \leq T} \rho^{T-t_{i}} U^{\text{source}}_{q_{i}}(g_{j}) = \sum_{i:t_{i} \leq T} \rho^{T-t_{i}} u^{\top}_{s_{i}} P_{s_{i}}$$
$$\mathbb{E}\left[U^{\text{source}}(G_{k},T+\delta)\right] = \mathbb{E}\left[U^{\text{source}}(G_{k'},T+\delta)\right]$$
$$u^{\top}_{s} P_{s} v = \mathbb{E}\left[U^{\text{source}}(G_{k'},T+\delta)\right] - \rho^{\delta} \mathbb{E}\left[U^{\text{source}}(G_{k'},T)\right]$$

since the ranking algorithms are by default trying to optimize the source side utility, it is unlikely that a group of source members would be consistently achieving poorer utility in a well-functioning system

ensuring that they receive comparable expected utility across groups over





Extension to Source Side Fairness ensuring that they receive comparable expected utility across groups over multiple sessions.

$$\begin{array}{ll} \text{Minimize } \sum_{k' \neq k} \left| u_s^\top P_s v - c_{k,k'} \right| \\ \text{Subject to } \mathbf{1}^\top P_s = 1, \quad P_s \mathbf{1} \leq 1, \quad \mathbf{0} \leq P_s(d,r) \leq 1, \quad f_{k,k'}^\top P_s w = \mathbf{0}, \quad \tilde{u}_s P_s w = \mathbf{0}. \end{array}$$

$$c_{k,k'} = \mathbb{E}[U^{\text{source}}(G_{k'}, T)]$$

since the ranking algorithms are by default trying to optimize the source side utility, it is unlikely that a group of source members would be consistently achieving poorer utility in a well-functioning system

$(+\delta)] - \rho^{\delta} \mathbb{E} \left[U^{\text{source}}(G_k, T) \right]$





Experiments (Simulation Setup)

- A graph to represent connections between members lacksquare
- Standard undirected graph, not bipartite
- A_t is the adjacency matrix at time t ullet

 $\operatorname{Aff}_{\operatorname{Graph}}(i,j) = A_t^2(i,j)$

$$\operatorname{Aff}_{\operatorname{member}}(i,j) = -\|X_i - X_j\|_2$$

 Graph Evolution A vertex i is picked at random. For each other vertex j, a "model" score s(i, j) is computed as

$$s(i,j) = p(G_i, G_j) \exp(\lambda \operatorname{Aff}_{\operatorname{Graph}}(i))$$



 $(i, j)) \exp(\mu \operatorname{Aff}_{\operatorname{member}}(i, j))$

Experiments

Fairness adjustments:

- 1. No fairness adjustment in the ranking (noReranker)
- 2. Equality of opportunity, using the primal for destination side fairness (primal)
- 3. Equality of opportunity, using the dual method for destination side fairness (dualNoDynamic)
- 4. Multi-Session destination utility adjustment via the dual (dualWithDynamic).

Experiments (Simulation Results)

Settings	Method	ΔDP	$\Delta \text{DP} $	Ratio of Source Utility	% Total Destination Utility
shown membe	rs noReranker	-0.0216	0.1964	0.5100	0.7051
m = 10	primal	0.0179	0.1671	0.5024	0.5371
(t) = 50	dualNoDynamic	0.0055	0.2125	0.5100	0.5860
rate of dual refresh	dualWithDynamic	0.0249	0.2008	0.5100	0.5706
m = 20 t = 50	noReranker	-0.0072	0.1182	0.5100	0.7059
	primal	0.0109	0.1002	0.5023	0.5269
	dualNoDynamic	0.0122	0.1090	0.5099	0.5536
	dualWithDynamic	0.0155	0.1099	0.5099	0.5371
m = 20 t = 20	noReranker	-0.0072	0.1182	0.5100	0.7059
	primal	0.0109	0.1002	0.5023	0.5269
	dualNoDynamic	0.0177	0.1105	0.5099	0.5210
	dualWithDynamic	0.0185	0.1090	0.5099	0.5166

Adding fairness re-rankers:

- help rebalance destination utility (last column)
- not significantly altering source utility metrics (see the 5th column)

Simulation Details: 1000 members and 1000 iterations

ast column) y metrics (see the 5th column)

Conclusion (My Thoughts)

- 1. Novel Idea by Designing multi-session dest-side utility Enables introducing multi-session dest-side utility and source-side utility
- 2. Using dual trick makes it flexible for large-scale implementation
- 3. Evaluation on only one fairness constraint (Demographic Parity) Since utility is available in test time Disparate Treatment could be evaluated too.
- 4. Does not provide different trade-offs between Utility and Fairness

5. Evaluation needed on Real-world Large-Scale data



